## TU-Chemnitz, Fakultät für Mathematik

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Homepage zur Übung: https://www.tu-chemnitz.de/mathematik/wire/WS1819/nla.php

## 7. Conjugate gradient method

1. Proof that: Let $V$ be a vector space with inner product (.,.), $P$ a subspace of $V$, and $f \in V$. The solution $\tilde{p}$ of the approximation problem

$$
\text { find } \tilde{p} \in P \text { s.t. }\|f-\tilde{p}\|=\min _{p \in P}\|f-p\|
$$

is defined by

$$
(f-\tilde{p}, p)=0 \quad \forall p \in P .
$$

[Hint: Assume that $(f-\tilde{p}, p)=a \neq 0$ and define a $q$ that gives a smaller norm, which is a contradiction.]
2. Implement the CG method and test it for the matrix $\mathrm{A}=$ gallery ('poisson', n ) for various $n$ and a random right hand side! To illustrate the performance you may plot the norms of the residual $\left\|r_{k}\right\|_{2}$ over the iteration number $k$.
3. For the cg method complete the proof from the lectures and show that

- $\left(r_{k+1}, p_{j}\right)=0 \quad \forall j=1,2, \ldots, k-1$.
- $\left(r_{k+1}, p_{j}\right)=\left(r_{k+1}, r_{j}\right) \quad \forall j=1,2, \ldots, k-1$.
- $\left(A p_{k+1}, p_{j}\right)=\left(r_{k+1}, r_{j}\right) \quad \forall j=1,2, \ldots, k-1$.

4. Proof that

$$
\beta_{k}=\frac{-\left(A r_{k+1}, p_{k}\right)}{\left(A p_{k}, p k\right)}=\frac{\left(r_{k+1}, r_{k+1}\right)}{\left(r_{k}, r_{k}\right)} .
$$

5. Proof that for the Chebyshev polynomials

$$
\tau_{0}(t)=1, \quad \tau_{1}(t)=t, \quad \tau_{k+1}(t)=2 t \tau_{k}(t)-\tau_{k-1}(t)
$$

the following idendity holds

$$
\tau_{k}(t)=\frac{1}{2}\left[\left(t+\sqrt{t^{2}-1}\right)^{k}+\left(t-\sqrt{t^{2}-1}\right)^{k}\right] .
$$

6. The following question is from An Introduction to the Conjugate Gradient Method Without the Agonizing Pain of J.R. Shewchuk:
Suppose you wish to solve $A x=b$ for a symmetric, positive-definite $N \times N$ matrix $A$. Unfortunately, the trauma of your linear algebra course has caused you to repress all memories of the Conjugate Gradient algorithm. Seeing you in distress, the Good Eigenfairy materializes and grants you a list of $d$ distinct eigenvalues (but not the eigenvectors) of $A$. However, you do not know how many times each eigenvalue is repeated.
Clever person that you are, you mumbled the following algorithm in your sleep this morning:
```
Choose an arbitrary starting point \(x_{0}\);
for \(i=0: d-1\)
    \(r_{i}=b-A * x_{i} ;\)
    Remove an arbitrary eigenvalue from the list and call it \(\lambda_{i}\);
    \(x_{i+1}=x_{i}+\lambda_{i}^{-1} r_{i} ;\)
end
```

No eigenvalue is used twice; on termination, the list is empty.
Show that upon termination of this algorithm, $x_{d}$ is the solution to $A x=b$.

